

An adjunction between categories of monads

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Abstract

Given a monad T on \mathcal{A} and a functor $G: \mathcal{A} \rightarrow \mathcal{B}$, if the right Kan extension of GT along G exists, then it has a canonical monad structure; we call this monad the pushforward of T along G and denote it by $G_{\#}T$. If F is left adjoint to G , then $G_{\#}T$ is the familiar monad structure on GTF . If T is the identity monad, then $G_{\#}T$ is the codensity monad of T . Even though it was introduced by Street [2] in the 1970s, the pushforward constructions has not yet been studied in depth. This talk will report on work towards this goal, the details of which can be found in [1].

Under suitable smallness and completeness conditions, $G_{\#}$ becomes a functor $\mathbf{Mnd}(\mathcal{A}) \rightarrow \mathbf{Mnd}(\mathcal{B})$. In fact, when G is full and faithful, $G_{\#}$ has a partial left adjoint $G^{\#}$, defined on those monads on \mathcal{B} whose endofunctor restricts to \mathcal{A} , i.e. those monads S such that there is an endofunctor S' of \mathcal{A} and an isomorphism $SG \cong GS'$. We denote the corresponding subcategory of $\mathbf{Mnd}(\mathcal{B})$ by $\mathbf{Mnd}(\mathcal{B})^{\text{res.}\mathcal{A}}$, giving an adjunction

$$\mathbf{Mnd}(\mathcal{A}) \begin{array}{c} \xrightarrow{G_{\#}} \\ \xleftarrow{G^{\#}} \end{array} \mathbf{Mnd}(\mathcal{B})^{\text{res.}\mathcal{A}} \subseteq \mathbf{Mnd}(\mathcal{B}),$$

with $G_{\#}$ full and faithful and $G^{\#}$ given by restriction.

The conditions which result in this situation are not hard to satisfy. For example, one may take G to be the inclusion $\mathbf{FinSet} \hookrightarrow \mathbf{Set}$, the inclusion $\mathbf{Set}_{\neq \emptyset} \hookrightarrow \mathbf{Set}$ (where the domain is the category of nonempty sets), the discrete-topology functor $\mathbf{Set} \rightarrow \mathbf{Top}$, or the forgetful functor $\mathbf{Field} \rightarrow \mathbf{CRing}$. The first two will be the guiding examples for this talk.

In the case of $\mathbf{Set}_{\neq \emptyset} \hookrightarrow \mathbf{Set}$, the category on the right-hand side of the adjunction is just $\mathbf{Mnd}(\mathbf{Set})$, and the corresponding reflective subcategory on the left is spanned by those monads all of whose pseudoconstants are induced by actual constants.

In the case of $\mathbf{FinSet} \hookrightarrow \mathbf{Set}$, the right-hand side of the adjunction includes monads such as the identity, the $(-) + E$ monad (for a finite set E), the $M \times (-)$ monad (for a finite monoid M), the powerset monad, the double-powerset monad, the ultrafilter monad, and the filter monad. I will explain what the pushforward of their restriction to \mathbf{FinSet} is; for example, starting with the powerset monad, this process produces the theory of continuous lattices. I will also explain why, with two trivial exceptions, pushforwards along $\mathbf{FinSet} \hookrightarrow \mathbf{Set}$ never have rank.

References

- [1] A. Doña Mateo, Pushforward monads, preprint arXiv:2406.15256, 2025
- [2] R. Street, *The formal theory of monads*, Journal of Pure and Applied Algebra 2.2 (1972)