

Lax-idempotent monads in homotopy theory

Joint with Fernando Abellán (work in progress)

Thomas Blom

Max Planck Institute for Mathematics, Bonn

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Example: Cocompletions of categories

- Let \mathcal{C} be a locally small $(\infty, 1)$ -category and $\mathcal{P}(\mathcal{C})$ its free cocompletion.
- Then \mathcal{C} is **cocomplete** if and only if $\mathcal{Y} : \mathcal{C} \rightarrow \mathcal{P}(\mathcal{C})$ admits a left adjoint L .

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- A functor $F: \mathcal{C} \rightarrow \mathcal{D}$ **preserves colimits** if and only if for any $X: I \rightarrow \mathcal{C}$, the assembly map

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- This is a **Beck–Chevalley condition**:

$$\begin{array}{ccc}
 \mathcal{P}(\mathcal{C}) & \xrightarrow{F_!} & \mathcal{P}(\mathcal{D}) \\
 \mathfrak{J} \uparrow & \circlearrowleft & \uparrow \mathfrak{J} \\
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 \end{array}
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This characterizes the sub- $(\infty, 2)$ -category $\mathcal{C}at_{(\infty, 1)}^{\text{cocomp}} \subseteq \mathcal{C}at_{(\infty, 1)}$:

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- Its **morphisms** are the maps $\mathcal{C} \rightarrow \mathcal{D}$ for which the Beck–Chevalley transformation

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is invertible.

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- A functor $p: \mathcal{C} \rightarrow \mathcal{B}$ between $(\infty, 1)$ -categories is a **cocartesian fibration** if and only if the inclusion

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- A functor $\mathcal{C} \rightarrow \mathcal{D}$ over \mathcal{B} preserves cocartesian morphisms if and only if

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- This completely characterizes $\mathrm{Cocart}(\mathcal{B}) \subseteq \mathrm{Cat}_{(\infty, 1)}/\mathcal{B}$.

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 - Their **objects** are those c for which $i: c \rightarrow Tc$ admits a left adjoint.
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These are examples of $(\infty, 2)$ -categories of algebras over a **lax-idempotent monad**.

Definition (after Kock¹ and Zöberlein²)

Let \mathcal{X} be an $(\infty, 2)$ -category and $T: \mathcal{X} \rightarrow \mathcal{X}$ a monad (i.e. an E_1 -algebra in $\text{End } \mathcal{X}$).

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Remark

Reversing the 1- or 2-morphisms gives three more variations:

- colax-idempotent monads: $\eta_{Tx} \dashv \mu_x$.
- lax-idempotent comonads: $\varepsilon_{Tx} \dashv \delta_x$.
- colax-idempotent comonads: $\delta_x \dashv \varepsilon_{Tx}$.

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4. The map $U: \text{Alg}(T) \rightarrow \mathcal{Y}$ is inverse to the canonical map $\mathcal{Y} \rightarrow \text{Alg}(T)$. □

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- One also obtains adjoint functor theorems for \mathcal{X} (cf. ADL³).

³N. Arkor, I. Di Liberti, and F. Loregian. “Adjoint functor theorems for lax-idempotent pseudomonads”. *Theory and Applications of Categories* 41.20 (2024), pp. 667–685.

Example: Cocompletions revisited

- Since $\mathcal{P}(\mathcal{C})$ is cocomplete, the **Yoneda embedding** $\mathcal{Y} : \mathcal{P}(\mathcal{C}) \rightarrow \mathcal{P}(\mathcal{P}(\mathcal{C}))$ always admits a left adjoint $L : \mathcal{P}(\mathcal{P}(\mathcal{C})) \rightarrow \mathcal{P}(\mathcal{C})$.

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- This defines the **multiplication map** of a lax-idempotent monad structure on $\mathcal{P} : \mathcal{Cat}_{(\infty,1)} \rightarrow \mathcal{Cat}_{(\infty,1)}$, with $\text{Alg}(\mathcal{P}) = \mathcal{Cat}_{(\infty,1)}^{\text{cocomp}}$.

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- This generalizes easily to the following two ways:
 - Given a collection of small $(\infty, 1)$ -categories \mathcal{K} , one can consider the free cocompletion $\mathcal{P}_{\mathcal{K}}(\mathcal{C})$ of \mathcal{C} with respect to **\mathcal{K} -indexed colimits**.
 - One can consider **enriched** or **internal** cocompletions.

Example: Operads

- Given a symmetric monoidal $(\infty, 1)$ -category \mathcal{C} , its **envelope** $\mathrm{Env}(\mathcal{C})$ is defined by the universal property

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- The resulting comonad on $\text{SymMon}_{(\infty, 1)}^{\text{strong}}$ has the property that the comultiplication $\delta_{\mathcal{C}}: \text{Env}(\mathcal{C}) \rightarrow \text{Env}(\text{Env}(\mathcal{C}))$ is **left adjoint** to the counit $\varepsilon_{\text{Env}(\mathcal{C})}: \text{Env}(\text{Env}(\mathcal{C})) \rightarrow \text{Env}(\mathcal{C})$.

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- Hence $\text{Env}: \text{SymMon}_{(\infty, 1)}^{\text{strong}} \rightarrow \text{SymMon}_{(\infty, 1)}^{\text{strong}}$ is a **colax-idempotent comonad**.
- Its coalgebras are **$(\infty, 1)$ -operads**:

$$\text{coAlg}(\text{Env}) \simeq \mathcal{Op}_{(\infty, 1)}.$$

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- The corresponding comonad Ind on $\text{Cat}_{(\infty,1)}^{\text{fil}}$ is **colax-idempotent**.
- \mathcal{C} is a coalgebra for this comonad if $\text{colim}: \text{Ind}(\mathcal{C}) \rightarrow \mathcal{C}$ admits a **further left adjoint**:

$$\begin{array}{ccc} & \widehat{\mathcal{J}} & \\ & \curvearrowright & \\ \mathcal{C} & \xleftarrow{\text{colim}} & \text{Ind}(\mathcal{C}) \\ & \curvearrowleft & \\ & \mathcal{J} & \end{array}$$

The diagram illustrates the relationship between a category \mathcal{C} and its Ind-construction $\text{Ind}(\mathcal{C})$. A horizontal arrow labeled colim points from $\text{Ind}(\mathcal{C})$ to \mathcal{C} . Above this arrow is a curved arrow pointing from \mathcal{C} to $\text{Ind}(\mathcal{C})$, labeled $\widehat{\mathcal{J}}$. Below the horizontal arrow is another curved arrow pointing from \mathcal{C} to $\text{Ind}(\mathcal{C})$, labeled \mathcal{J} . Both curved arrows have a \perp symbol positioned between them.

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The diagram illustrates the relationship between a category \mathcal{C} and its ind-completion $\text{Ind}(\mathcal{C})$. A central horizontal arrow points from $\text{Ind}(\mathcal{C})$ to \mathcal{C} and is labeled colim . Above this arrow is a curved arrow pointing from \mathcal{C} to $\text{Ind}(\mathcal{C})$, labeled $\widehat{\mathcal{C}}$. Below the central arrow is another curved arrow pointing from \mathcal{C} to $\text{Ind}(\mathcal{C})$, labeled \perp . The entire diagram is enclosed in a large, faint, light-blue arrow pointing from left to right.

- If such a \mathcal{C} is furthermore accessible, then it is called **compactly assembled**.

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- If such a \mathcal{C} is furthermore accessible, then it is called **compactly assembled**.
- Efimov⁴ showed that the usual K -theory functor extends uniquely to a **localizing invariant** $K^{\text{cts}}: \text{CompAss}_{\text{Stable}} \rightarrow \text{Sp}$.

⁴A. I. Efimov. *K-theory and localizing invariants of large categories*. 2025. arXiv: 2405.12169.

What determines a lax-idempotent monad?

- To any lax idempotent monad T on an $(\infty, 2)$ -category \mathcal{X} , one can associate the collection of **unit maps**

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- From this collection one can recover the (locally full) $(\infty, 2)$ -category $\mathrm{Alg}(T) \subset \mathcal{X}$:
 - Its **objects** are those x for which η_x admits a left adjoint α_x .
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- This suggests a 1-1 correspondence between certain collections of arrows in \mathcal{X} and lax-idempotent monads on \mathcal{X} .

A non-coherent characterization of lax-idempotent monads

Theorem (after Marmolejo–Wood⁵)

Let $\mathcal{U} = \{x \multimap y\} \subseteq \text{Ar}(\mathcal{X})^\simeq$ be a collection of arrows such that:

1. $\text{ev}_0: \mathcal{U} \rightarrow \mathcal{X}$ is essentially surjective.

⁵F. Marmolejo and R. J. Wood. “Kan extensions and lax idempotent pseudomonads”. *Theory and*

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$$\begin{array}{ccccc} x & & x' & \xrightarrow{g} & z \\ \downarrow \circlearrowleft & \searrow f & \downarrow \circlearrowleft & \nearrow \text{Lan } g & \\ y & \xrightarrow{\text{Lan } f} & y' & & \end{array}$$
 exhibits $(\text{Lan } g) \circ (\text{Lan } f)$ as a left Kan extension of $(\text{Lan } g) \circ f$.

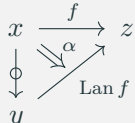
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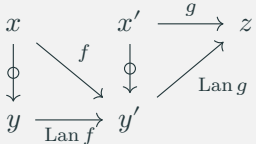
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Then there exists a lax-idempotent monad on \mathcal{X} given by “ $x \mapsto \text{target}(x \multimap y)$ ”, and any lax-idempotent monad arises on \mathcal{X} this way.

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Proof sketch.

Let $j: x \multimap y$ in \mathcal{U} and $z \in \text{ev}_1(\mathcal{U})$ be given.

Call $y \rightarrow z$ **extended** if it is left Kan extended along $j: x \multimap y$. (This turns out to depend only on y and not on j .)

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We define $\mathcal{Y} \subseteq \mathcal{X}$ to be the locally full subcategory with

- **Objects:** $\{z \mid \text{there exists } x \multimap y \text{ in } \mathcal{U} \text{ with } y \simeq z\}$.
- **Morphisms:** $\{f: z \rightarrow z' \mid f \text{ is extended}\}$.

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The corresponding monad is **lax-idempotent**. □

 Thank you for listening! 