

Double-category of sites

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Introduction

(Grothendieck) topoi can be presented through sites.

As well, geometric morphisms can be presented either:

in a contravariant way, by:

morphisms of sites

characterized by some
cover-preserving property

in a covariant way, by:

comorphisms of sites

characterized by some
cover-lifting property

Can those two classes be mixed in a same 2-categorical structure on sites ?

Problem: morphisms and comorphisms of sites do not compose with each other !

Introduction

Solution: make them the horizontal and vertical arrows of a *double-category* of sites.

Such a structure does not require them to compose with each other.

Moreover the sheaf topos construction arranges nicely into a double-functor.

Some results of topos theory will rephrase as double-categorical statements.

This talk is based on the preprint:

O. Caramello and A. Osmond. "Morphisms and comorphisms of sites I – Double categories of sites". In: *arXiv:2505.08766* (2025)

Outlay

Morphisms and comorphisms of sites

Double-category of sites and sheafification double-functor

(Co)morphisms as (co)lax morphisms of coalgebras

Morphisms and comorphisms of sites

Definition

A *sieve* on an object c in a category \mathcal{C} is a subobject of the representable $S \hookrightarrow \mathcal{Y}_c$.

A *coverage* J on \mathcal{C} consists for each c of a set $J(c)$ of sieves on c such that

- for each c , the *maximal sieve*, which is \mathcal{Y}_c , is in $J(c)$
- for each arrow $a : d \rightarrow c$ and each S in $J(c)$, the *pullback sieve* below is in $J(d)$

$$a^*S = \left\{ v : d' \rightarrow d \mid \exists u : c' \rightarrow c \in S \text{ and a factorization } \begin{array}{ccc} d' & \xrightarrow{\exists} & c' \\ v \downarrow & & \downarrow u \in S \\ d & \xrightarrow{a} & c \end{array} \right\}$$

We will assume coverages to be *sifted*: if $S \leq R$ and $S \in J(c)$, then $R \in J(c)$.

A *site* is a pair (\mathcal{C}, J) with J a (sifted) coverage on \mathcal{C} .

Any (small generated) site (\mathcal{C}, J) induces a topos $\mathbf{Sh}(\mathcal{C}, J)$.

Notions of morphisms between sites

A functor between sites may behave in two relevant ways relative to the coverages:

- either by *preserving* covering sieves;
- either by *lifting* covering sieves.

Combined with flatness, cover-preservation defines a notion of *morphism of sites*.

On the other hand cover-lifting defines a notion of *comorphism of sites*.

Both induce geometric morphisms between associated sheaf topoi.

Let us revisit those ideas through the formalism of *extension and restriction*.

Extension and restriction

Recall that any functor $f : \mathcal{C} \rightarrow \mathcal{D}$ induces a triple of adjoints

$$\begin{array}{ccc} & \text{lex}_f & \\ \widehat{\mathcal{C}} & \begin{array}{c} \xrightarrow{\perp} \\ \text{rest}_f \\ \xleftarrow{\perp} \end{array} & \widehat{\mathcal{D}} \\ & \text{rext}_f & \end{array}$$

where lex_f (resp. rext_f) sends a presheaf $X : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}$ to

$$\text{lex}_f X = \text{lan}_{f^{\text{op}}} X \qquad \text{rext}_f X = \text{ran}_{f^{\text{op}}} X$$

while rest_f sends a presheaf $Y : \mathcal{D}^{\text{op}} \rightarrow \mathbf{Set}$ to

$$\text{rest}_f Y = Y \circ f^{\text{op}}$$

Beware that lex and rext are covariant in f , while rest is contravariant in f .

Extension and restriction of sieves

For a sieve $S \rhd \mathcal{J}_c$ in \mathcal{C} , the left extension $\text{lex}_f S$ induces a sieve on $f(c)$ by taking the image

$$\begin{array}{ccc} \text{lex}_f(S) & \xrightarrow{\text{lex}_f(m)} & \mathcal{J}_{f(c)} \\ & \searrow & \nearrow f[m] \\ & f[S] & \end{array}$$

corresponding to the set of arrows

$$\left\{ v : d \rightarrow f(c) \mid \begin{array}{ccc} d & \xrightarrow{v} & f(c) \\ \exists \downarrow & \nearrow f(u), u \in S & \\ f(c') & & \end{array} \right\}$$

For a sieve $R \rhd \mathcal{J}_d$ in \mathcal{D} , the restriction $\text{res}_f(R)$ induces a sieve on c by taking the pullback

$$\begin{array}{ccc} f^{-1}(R) & \xrightarrow{\quad} & \text{res}_f R \\ \downarrow & \lrcorner & \downarrow \\ \mathcal{J}_c & \xrightarrow{1_{f(c)}} & \mathcal{D}(f, f(c)) \end{array}$$

corresponding to the set of arrows

$$\left\{ u : c' \rightarrow c \mid \begin{array}{ccc} f(c') & \xrightarrow{f(u)} & f(c) \\ \exists \downarrow & \nearrow v \in R & \\ d' & & \end{array} \right\}$$

Morphisms and comorphisms of sites

Let $f : (\mathcal{C}, J) \rightarrow (\mathcal{D}, K)$ be a functor between sites; then :

Definition

f is *cover-preserving* if for any c in \mathcal{C} and any $S \rightrightarrows \mathcal{J}_c$ in $J(c)$, the sieve $f[S]$ is in $K(f(c))$.

A *morphism of sites* is a **flat** functor that is cover-preserving.

Call \mathbf{Sit}^b the 2-category of

- sites
- morphisms of sites
- and transformations

Definition

f is *cover-lifting* if for any c in \mathcal{C} and any $R \rightrightarrows \mathcal{J}_{f(c)}$ in $K(f(c))$, the sieve $f^{-1}(R)$ is in $J(c)$.

A *comorphism of sites* is a functor that is cover-lifting.

Call \mathbf{Sit}^\sharp the 2-category of

- sites
- comorphisms of sites
- and transformations

Geometric morphisms induced from (co)morphisms of sites

A morphism of sites

$$f : (\mathcal{C}, J) \rightarrow (\mathcal{D}, K)$$

induces a geometric morphism

$$\mathbf{Sh}(f) : \widehat{\mathcal{D}}_K \rightarrow \widehat{\mathcal{C}}_J$$

with inverse image

$$\begin{array}{ccc} \widehat{\mathcal{C}} & \xrightarrow{\text{lex}_f} & \widehat{\mathcal{D}} \\ i_J \uparrow & & \downarrow a_K \\ \widehat{\mathcal{C}}_J & \xrightarrow{\mathbf{Sh}(f)^*} & \widehat{\mathcal{D}}_K \end{array}$$

This defines a 1-contravariant,
2-covariant pseudofunctor

$$(\mathbf{Sit}^b)^{\text{op}} \xrightarrow{\mathbf{Sh}} \mathbf{Top}$$

A comorphism of sites

$$F : (\mathcal{C}, J) \rightarrow (\mathcal{D}, K)$$

induces a geometric morphism

$$C_F : \widehat{\mathcal{C}}_J \rightarrow \widehat{\mathcal{D}}_K$$

with inverse image

$$\begin{array}{ccc} \widehat{\mathcal{D}} & \xrightarrow{\text{rest}_F} & \widehat{\mathcal{C}} \\ i_J \uparrow & & \downarrow a_K \\ \widehat{\mathcal{D}}_L & \xrightarrow{C_F^*} & \widehat{\mathcal{C}}_J \end{array}$$

This defines a 1-covariant,
2-contravariant pseudofunctor

$$(\mathbf{Sit}^\sharp)^{\text{co}} \xrightarrow{C} \mathbf{Top}$$

Double-category of sites and sheafification double-functor

Double-category of sites

Definition

We define the double-category **Sit**^h as having

- as objects (small generated) sites,
- as horizontal arrows morphisms of sites,
- as vertical arrows comorphisms of sites,
- and as a double-cells lax squares as below, with $\begin{cases} f, h \text{ morphisms of sites} \\ G, K \text{ comorphisms of sites} \end{cases}$

$$\begin{array}{ccc} (\mathcal{A}, M) & \xrightarrow{f} & (\mathcal{B}, L) \\ G \downarrow & \phi & \downarrow K \\ (\mathcal{C}, J) & \xrightarrow{h} & (\mathcal{D}, K) \end{array}$$

$$\begin{array}{ccc} (\mathcal{A}, M) & \xrightarrow{f} & (\mathcal{B}, L) \\ G \downarrow & \swarrow \phi & \downarrow K \\ (\mathcal{C}, J) & \xrightarrow{h} & (\mathcal{D}, K) \end{array}$$

Constructing a sheafification double-functor

We want a double-functor $\left\{ \begin{array}{l} \text{with horizontal component the pseudofunctor } \mathbf{Sh} \\ \text{with vertical component the pseudofunctor } \mathbf{C} \end{array} \right.$

For double-cell, suppose one has a 2-cell of the following form:

$$\begin{array}{ccc} \mathcal{A} & \xrightarrow{f} & \mathcal{B} \\ g \downarrow & \phi \swarrow & \downarrow k \\ \mathcal{C} & \xrightarrow{h} & \mathcal{D} \end{array}$$

Such a square induces a cross-adjoint square constructed from the composite 2-cell

$$\begin{array}{ccc} \hat{\mathcal{A}} & \xrightarrow{\text{lex}_f} & \hat{\mathcal{B}} \\ \text{rest}_g \uparrow & \overline{\phi} \swarrow & \uparrow \text{rest}_k \\ \hat{\mathcal{C}} & \xrightarrow{\text{lex}_h} & \hat{\mathcal{D}} \end{array}$$

$$\begin{array}{ccc} \text{lex}_f \text{rest}_g & \xRightarrow{\eta_h} & \text{lex}_f \text{rest}_g \text{rest}_h \text{lex}_h \\ \overline{\phi} \Downarrow & & \Downarrow \text{rest}_\phi \\ \text{rest}_k \text{lex}_h & \xleftarrow{\epsilon_f} & \text{lex}_f \text{rest}_f \text{rest}_k \text{lex}_h \end{array}$$

Sheafification of double-cells

If now one has sites (\mathcal{A}, M) , (\mathcal{B}, L) , (\mathcal{C}, J) and (\mathcal{D}, K) , related through a lax square

$$\begin{array}{ccc} (\mathcal{A}, M) & \xrightarrow{f} & (\mathcal{B}, L) \\ G \downarrow & \swarrow \phi & \downarrow K \\ (\mathcal{C}, J) & \xrightarrow{h} & (\mathcal{D}, K) \end{array}$$

f, h morphisms of sites
 G, K comorphisms of sites

then the sheafification functor $\alpha_L : \widehat{\mathcal{B}} \rightarrow \widehat{\mathcal{B}}_L$ sends $\overline{\phi}$ a 2-cell ϕ^b corresponding to the inverse image part of a geometric transformation

$$\begin{array}{ccc} \widehat{\mathcal{A}}_M & \xrightarrow{\widehat{f}^*} & \widehat{\mathcal{B}}_L \\ c_G^* \uparrow & \searrow \phi^b & \uparrow c_K^* \\ \widehat{\mathcal{C}}_J & \xrightarrow{\widehat{h}^*} & \widehat{\mathcal{D}}_K \end{array} \qquad \begin{array}{ccc} \widehat{\mathcal{A}}_M & \xleftarrow{\text{Sh}(f)} & \widehat{\mathcal{B}}_L \\ c_G \downarrow & \searrow \widehat{\phi} & \downarrow c_K \\ \widehat{\mathcal{C}}_J & \xleftarrow{\text{Sh}(h)} & \widehat{\mathcal{D}}_K \end{array}$$

This is a 2-cell in the *lax quintet* double-category **Top**[□] of Grothendieck topoi.

Sheafification as a double-functor

Theorem

Sheafification defines a surjective on object double-functor into the lax quintet double-category of topoi , which is

- *horizontally contravariant with horizontal component \mathbf{Sh}*
- *vertically covariant with vertical component \mathbf{C}*

$$\begin{array}{ccccc} (\mathbf{Sit}^b)^{\text{op}} & \xrightarrow{h} & (\mathbf{Sit}^{\sharp})^{\text{hop}}_{\text{vco}} & \xleftarrow{v} & (\mathbf{Sit}^{\sharp})^{\text{co}} \\ & \searrow \mathbf{Sh} & \downarrow & \swarrow \mathbf{C} & \\ & & \mathbf{Top}^{\square} & & \end{array}$$

Lax squares inverted by this double-functor generalize the notion of *exact squares*.

Exact squares

A lax square in **Cat** is *exact* if the associated extension/restriction 2-cell is invertible

$$\begin{array}{ccc}
 \mathcal{A} & \xrightarrow{f} & \mathcal{B} \\
 g \downarrow & \phi \swarrow & \downarrow k \\
 \mathcal{C} & \xrightarrow{h} & \mathcal{D}
 \end{array}
 \qquad
 \begin{array}{ccc}
 \widehat{\mathcal{A}} & \xrightarrow{\text{lex}_f} & \widehat{\mathcal{B}} \\
 \text{rest}_g \uparrow & \bar{\phi} \swarrow \simeq & \uparrow \text{rest}_k \\
 \widehat{\mathcal{C}} & \xrightarrow{\text{lex}_h} & \widehat{\mathcal{D}}
 \end{array}$$

Definition

A lax square as below left (underlying a double-cell of **Sit**^h) will be said *locally exact* if the corresponding transformation below right is invertible

$$\begin{array}{ccc}
 (\mathcal{A}, M) & \xrightarrow{f} & (\mathcal{B}, L) \\
 G \downarrow & \phi \swarrow & \downarrow K \\
 (\mathcal{C}, J) & \xrightarrow{h} & (\mathcal{D}, K)
 \end{array}
 \qquad
 \begin{array}{ccc}
 \widehat{\mathcal{A}}_M & \xrightarrow{\text{Sh}(f)^*} & \widehat{\mathcal{B}}_L \\
 c_G^* \uparrow & \bar{\phi}^b \swarrow & \uparrow c_K^* \\
 \widehat{\mathcal{C}}_J & \xrightarrow{\text{Sh}(h)^*} & \widehat{\mathcal{D}}_K
 \end{array}$$

This admits a concrete characterization in term of *relative cofinality* à la [2].

Comma and cocomma squares in \mathbf{Sit}^{\natural}

In \mathbf{Cat} the ur-examples of exact squares are comma and cocomma; similarly:

Proposition

If $\begin{cases} f: (\mathcal{C}, J) \rightarrow (\mathcal{D}, K) \text{ morph.} \\ G: (\mathcal{B}, L) \rightarrow (\mathcal{D}, K) \text{ comorph.} \end{cases}$

then there is a topology $J_{G,f}$ on $G \downarrow f$ that makes the comma square a double-cell of \mathbf{Sit}^{\natural}

$$\begin{array}{ccc} (G \downarrow f, J_{G,f}) & \xrightarrow{\pi_0} & (\mathcal{B}, L) \\ \pi_1 \downarrow & \lambda_{G,f} & \downarrow G \\ (\mathcal{C}, J) & \xrightarrow{f} & (\mathcal{D}, K) \end{array}$$

Proposition

If $\begin{cases} f: (\mathcal{C}, J) \rightarrow (\mathcal{B}, L) \text{ morph.} \\ G: (\mathcal{C}, J) \rightarrow (\mathcal{D}, K) \text{ comorph.} \end{cases}$

then there is a topology $J_{f,G}$ on $f \uparrow G$ that makes the cocomma square a double-cell of \mathbf{Sit}^{\natural}

$$\begin{array}{ccc} (\mathcal{C}, J) & \xrightarrow{f} & (\mathcal{B}, L) \\ G \downarrow & \lambda_{f,G} & \downarrow \iota_0 \\ (\mathcal{D}, K) & \xrightarrow{\iota_1} & (f \uparrow G, J_{f,G}) \end{array}$$

In both cases, the underlying square in \mathbf{Cat} is exact.

Some effects of the mixed variance of **Sh**

Vertical and horizontal cells can be related either through:

Definition

- *conjoint squares*

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \parallel & \epsilon & \downarrow G \\ A & = & A \end{array} \quad \begin{array}{ccc} B & = & B \\ G \downarrow & \eta & \parallel \\ A & \xrightarrow{f} & B \end{array}$$

$$\epsilon \bullet \eta = 1_G \quad \eta \circ \epsilon = \text{id}_f$$

Definition

- *companion squares*

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ G \downarrow & \phi & \parallel \\ B & = & B \end{array} \quad \begin{array}{ccc} A & = & A \\ \parallel & \psi & \downarrow G \\ A & \xrightarrow{f} & B \end{array}$$

$$\phi \bullet \psi = 1_G \quad \psi \circ \phi = \text{id}_f$$

By mixed variance, **Sh** sends $\begin{cases} \text{companions to conjoints} \\ \text{conjoints to companions} \end{cases}$

Those squares are alike those in **Cat** whose exactness captures:

- the fact of being adjoint
- the fact of being respectively fully faithful and absolutely dense.

Adjunction through exactness

A functor f is right adjoint to G in **Cat** iff there is an exact square as below

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{f} & \mathcal{D} \\ \parallel & \xleftarrow{\epsilon} & \downarrow G \\ \mathcal{C} & \xlongequal{\quad} & \mathcal{C} \end{array} \quad \text{lex}_f \simeq \text{rest}_G$$

Definition

A morphism of site $f : (\mathcal{C}, J) \rightarrow (\mathcal{D}, K)$ will be said to be *weakly right adjoint* to a comorphism G if there exist a locally exact square as below:

$$\begin{array}{ccc} (\mathcal{C}, J) & \xrightarrow{f} & (\mathcal{D}, K) \\ \parallel & \xleftarrow{\epsilon} & \downarrow G \\ (\mathcal{C}, J) & \xlongequal{\quad} & (\mathcal{C}, J) \end{array}$$

Proposition

f is weakly right adjoint to G iff f and G induce the same geometric morphism

$$\mathbf{Sh}(f) \simeq C_G : \mathbf{Sh}(\mathcal{D}, K) \rightarrow \mathbf{Sh}(\mathcal{C}, J)$$

Local exactness criterion for local and totally connected morphisms

A functor $f : \mathcal{C} \rightarrow \mathcal{D}$ is fully faithful iff its identity square below is exact:

$$\begin{array}{ccc} \mathcal{C} & \xlongequal{\quad} & \mathcal{C} \\ \parallel & & \downarrow f \\ \mathcal{C} & \xrightarrow{f} & \mathcal{D} \end{array} \quad \text{rest}_f \text{lex}_f = \text{id}$$

A functor $f : \mathcal{C} \rightarrow \mathcal{D}$ is *absolutely dense* iff the square below is exact:

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{f} & \mathcal{D} \\ f \downarrow & & \parallel \\ \mathcal{D} & \xlongequal{\quad} & \mathcal{D} \end{array} \quad \text{lex}_f \text{rest}_f = 1$$

Proposition

Let f be both $\begin{cases} \text{a morphism} \\ \text{a comorphism} \end{cases}$

$\text{Sh}(f)$ is a local geometric morphism iff this square is locally exact:

$$\begin{array}{ccc} (\mathcal{C}, J) & \xlongequal{\quad} & (\mathcal{C}, J) \\ \parallel & & \downarrow f \\ (\mathcal{C}, J) & \xrightarrow{f} & (\mathcal{D}, K) \end{array}$$

Proposition

Let f be both $\begin{cases} \text{a morphism} \\ \text{a comorphism} \end{cases}$

C_f is a totally connected geometric morphism iff this square is locally exact:

$$\begin{array}{ccc} (\mathcal{C}, J) & \xrightarrow{f} & (\mathcal{D}, K) \\ f \downarrow & & \parallel \\ (\mathcal{D}, K) & \xlongequal{\quad} & (\mathcal{D}, K) \end{array}$$

(Co)morphisms as (co)lax morphisms of coalgebras

Why a double-categorical structure ?

Why morphisms and comorphisms do arrange in a double-category ?

This double-category is an instance of a special family of double-categories.

double-categories of profunctors have a different flavour and are less symmetrical.

Here horizontal and vertical maps are two classes of functors with dual properties, rather than a class of functors and a class of relations generalizing them.

This reminds a more symmetric kind of double-categories, those that arise as double-categories of (co)algebras with lax and colax morphisms for (co)monads !

Coalgebra and (co)lax morphisms for a copointed endofunctor

Definition

Let \mathcal{K} be a 2-category and $T : \mathcal{K} \rightarrow \mathcal{K}$ a copointed endo-2-functor;

A *coalgebra* for (T, ε) is a pair (C, γ) with C in \mathcal{K} and α a section of the counit

$$\begin{array}{ccc} C & \xrightarrow{\gamma} & TC \\ & \searrow & \downarrow \varepsilon_C \\ & & C \end{array}$$

A *lax* (resp. *colax*) *morphism* of coalgebras $(C, \gamma) \rightarrow (D, \delta)$ is a pair (f, ϕ) with $f : C \rightarrow D$ in \mathcal{K} and ϕ a 2-cell as on the left (resp. on the right)

$$\begin{array}{ccc} C & \xrightarrow{f} & D \\ \gamma \downarrow & \xleftarrow{\phi} & \downarrow \delta \\ TC & \xrightarrow{Tf} & TD \end{array}$$

$$\begin{array}{ccc} C & \xrightarrow{f} & D \\ \gamma \downarrow & \xrightarrow{\phi} & \downarrow \delta \\ TC & \xrightarrow{Tf} & TD \end{array}$$

whose pasting along the naturality square of the pointer is an identity 2-cell.

Double-category of coalgebras, lax and colax morphisms

Proposition (After Paré-Grandis [3])

For any copointed endo-2-functor T , one can form a double-category $T\text{-coAlg}$ of strict coalgebras, lax morphisms as horizontal cells, colax morphisms as vertical cells, and as 2-cell, the lax squares of the form

$$\begin{array}{ccc} (A, \alpha) & \xrightarrow{(f, \phi)} & (B, \beta) \\ (h, \eta) \downarrow & \xRightarrow{\psi} & \downarrow (k, \chi) \\ (C, \gamma) & \xrightarrow{(g, \kappa)} & (D, \delta) \end{array}$$

The double-cells of this double-category consist hence in 2-cells $\psi : gh \Rightarrow kf$ intertwining the lax and colax morphism structures in the following coherence

$$\begin{array}{c} \begin{array}{ccccc} & & B & \xrightarrow{\beta} & TB \\ & f \nearrow & \leftarrow \phi = Tf & \nearrow & \\ A & \xrightarrow{\alpha} & TA & & \\ & h \searrow & \xRightarrow{\eta} Th & \xRightarrow{T\psi} & TD \\ & & C & \xrightarrow{\gamma} & TC \end{array} \\ \end{array} = \begin{array}{c} \begin{array}{ccccc} & & B & \xrightarrow{\beta} & TB \\ & f \nearrow & & \searrow k & \\ A & & & & D \xrightarrow{\delta} TD \\ & h \searrow & \xRightarrow{\psi} & \nearrow g & \\ & & C & \xrightarrow{\gamma} & TC \end{array} \end{array}$$

The category $\mathbb{S}\mathcal{C}$

Definition

Define for each category \mathcal{C} the category $\mathbb{S}\mathcal{C}$ as having:

- as objects pairs (c, F) with c an object of \mathcal{C} and F a filter of $\text{Sub}_{\widehat{\mathcal{C}}} \downarrow_c$
- as morphisms $(c, F) \rightarrow (c', F')$ morphisms $u : c \rightarrow c'$ such that

$$F' \leq (u^*)^{-1}(F)$$

This category is fibered over \mathcal{C} with posetal fibers $\mathbb{F}_{\mathcal{C}}(c)$ at each c

$$\mathbb{S}\mathcal{C} \xrightarrow{\pi_{\mathcal{C}}} \mathcal{C}$$

For a functor $f : \mathcal{C} \rightarrow \mathcal{D}$, one can define at each filter F of $\text{Sub}_{\widehat{\mathcal{C}}} \downarrow_c$ the filter $f[F]$ generated from the set of sieves of the form $f[S]$ for S in F .

Defines the functor $\mathbb{S}f : \mathbb{S}\mathcal{C} \rightarrow \mathbb{S}\mathcal{D}$ sending (c, F) to the pair $(f(c), f[F])$.

The copointed endo-2-functor \mathbb{S}

Hence this construction is functorial on **Cat**: we have an endo-2-functor

$$\mathbf{Cat} \xrightarrow{\mathbb{S}} \mathbf{Cat}$$

Moreover this endofunctor is copointed through the projections $\pi_{\mathcal{C}} : \mathbb{S}\mathcal{C} \rightarrow \mathcal{C}$, whose naturality produce morphisms of fibrations

$$\begin{array}{ccc} \mathbb{S}\mathcal{C} & \xrightarrow{\mathbb{S}f} & \mathbb{S}\mathcal{D} \\ \pi_{\mathcal{C}} \downarrow & & \downarrow \pi_{\mathcal{D}} \\ \mathcal{C} & \xrightarrow{f} & \mathcal{D} \end{array}$$

Proposition

A coalgebra structure on \mathcal{C} for the copointed endofunctor (\mathbb{S}, π) is a coverage on \mathcal{C} .

Indeed a section $J : \mathcal{C} \rightarrow \mathbb{S}\mathcal{C}$ of $\pi_{\mathcal{C}}$ picks for each c a filter of sieves on c .

Moreover functoriality says that for any $u : c \rightarrow c'$, one has a restriction

$$\begin{array}{ccc} J(c) & \hookrightarrow & \text{Sub}_{\widehat{\mathcal{C}}} \downarrow_c \\ \uparrow & & \uparrow u^* \\ J(c') & \hookrightarrow & \text{Sub}_{\widehat{\mathcal{C}}} \downarrow_{c'} \end{array}$$

expressing that for any R in $J(c')$ the pullback sieve u^*R is in $J(c)$.

This is exactly what a coverage is!

(Co)morphisms of sites as (co)lax morphisms of coalgebras

Proposition

A functor is a lax morphism of coalgebras iff it is cover-preserving.

Having a 2-cell as below

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{f} & \mathcal{D} \\ J \downarrow & \xleftarrow{\phi} & \downarrow K \\ \mathbb{S}\mathcal{C} & \xrightarrow{\mathbb{S}f} & \mathbb{S}\mathcal{D} \end{array}$$

amounts to an inequality at each c

$$f[J(c)] \leq K(f(c))$$

which means that for any S in $J(c)$, $f[S]$ is in $K(f(c))$.

Proposition

A functor is a colax morphism of coalgebras iff it is cover-lifting.

Having a 2-cell as below

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{f} & \mathcal{D} \\ J \downarrow & \xrightarrow{\phi} & \downarrow K \\ \mathbb{S}\mathcal{C} & \xrightarrow{\mathbb{S}f} & \mathbb{S}\mathcal{D} \end{array}$$

amounts to an inequality at each c

$$K(f(c)) \leq f[J(c)]$$

which means that any R in $K(f(c))$ contains some $f[S]$ for a S in $J(c)$.

Further aspects and future directions

There are still several questions regarding this double-categorical approach:

- actually \mathbb{S} bears a structure of 2-comonad, but sites are only *normal lax* coalgebras; it is still unclear what strict coalgebras for the full comonad are.
- this can be fixed by an indexed form of this comonad, where sites correctly corresponds to coalgebras;
- combine this with flatness s morphisms of sites are lax morphisms of coalgebras ?
- in [4] **Top** was shown to be a bilocalization of \mathbf{Sit}^b at *dense morphisms of sites*. Is similarly \mathbf{Top}^\square a double-localization of \mathbf{Sit}^b ?
- in a upcoming work we will also show how morphisms and comorphisms are subsumed by *continuous distributors* between sites;
- is \mathbf{Sit}^b a good framework to do some formal category theory sites mixing the Yoneda structure of **Cat** with coverages ?

Thank you for your attention !

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